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# Towards Generating Robust Key Based on Neural Networks and Chaos Theory 

Abeer Tariq Maolood*, Ala'a Talib Khudhair<br>Department of Computer Science, University of Technology, Baghdad, Iraq


#### Abstract

There are large numbers of weakness in the generated keys of security algorithms. This paper includes a new algorithm to generate key of 5120 bits for a new proposed cryptography algorithm for 10 rounds that combine neural networks and chaos theory (1D logistic map). Two methods of neural networks (NN) are employed as Adaline and Hopfield and the results are combined through several sequential operation. Carefully integrating high quality random number generators from neural networks and chaos theory to obtain suitable key for randomness and complexity.


Keywords: key generation, neural network, Adaline, Hopfield, chaos theory


الخلاصة


## 1. Introduction

Cryptography is a technique which converts a given message into an unreadable format which can be seen as a noise by the third parties. It is a vital part of a safe communication process .Cryptography is a skill of sending the data in such a form that only those for whom it is intended can read it. Any cryptography algorithm needs to generate key that can be used in the steps of cryptography [1]. In this paper a new technique is employed for key generation by using Artificial Neural Network (ANN) as Hopfield and Adaline. Hopfield invented by the physician (John Hopfield) in 1982 which has these features: unsupervised learning, associative memory, full connection, single layer, feedback, fixed weight, bipolar, linear activation function and input nodes equal to output nodes. While Adaline NN invented by psychologist Donald Hebb postulated in 1949 [1]. Generating random numbers are mainly used to generate secret keys or random sequences and it can be carried out by many different techniques such as chaotic maps. In recent years many pseudo random numbers or sequences generators are proposed based on chaotic maps. The most important thing in any cryptography algorithms is having secure key and there are several requirements for achieving secure key. One of these requirements is randomize and key length, etc. [2].

[^0]
## Related work

This section discusses important previous works of Hopfield, Adaline and chaos:

1. Marat Akhmet \& Mehmet Onur Fen, 2014, has applied a Generation of cyclic/toroidal chaos by Hopfield neural networks. Have proposed system to discuss the appearance of cyclic and toroidal chaos in Hopfield neural networks [3].
2. Terence Kwok \& Kate Smith\& Lipo Wang, 2013, Incorporating Chaos into the Hopfield Neural network for Combinatorial Optimization, the proposed method consisted of adding chaotic noise to a Hopfield network. The N -Queen problem is used as an application to test and compare the performance and robustness of the three methods. The traditional simulated annealing is also included for comparison in order to contrast the effectiveness of the various approaches [4].
3. Delhi Technological \& Swati Mishra, 2013, Public key cryptography produced from neural networks and genetic algorithms. This paper made use of Hebbian learning rule to train the ANN of both sender and receiver machines. In the field of Public Key Cryptography (PKC), Pseudo Random Number Generator (PRNG) are widely used to generate unique keys and random numbers used in ANN which are found to possess many types of possible attacks. It is essential for a key to possess randomness for key strength and security. The PKC was generated by using of ANN through Genetic Algorithm (GA) [5].

## 2. Theoretical background

### 2.1 Hopfield NN

Hopfield Neural Network is an example of the network with feedback (so-called recurrent network), where outputs of neurons are connected to input of every neuron by means of the appropriate weights. Of course there are also inputs which provide neurons with components of test vector.[6]


Figure 1- Hopfield NN Topology

## Hopfield NN Learning

Initialize ( N ) no. of input node.

$$
F(x)= \begin{cases}1 & \text { if } x>0 \\ -1 & \text { if } x \leq 0\end{cases}
$$

-Convert 0 to -1
-Compute weight matrix.
$W_{i, j}=\left\{\begin{array}{cc}0 & \text { if } i=j \\ \sum_{i=1}^{p} x_{i} \times x_{j} & \text { otherwise }\end{array}\right.$ where $\mathrm{i}, \mathrm{j}=1.2 \ldots . \mathrm{N}[6]$


```
    Wait \((i, j)=\sum x_{i} x_{j}\)
Step4:final wait(i, j)
Next j
Next i
End
```


### 2.2 Adaline neural network

Adaline which stands for Adaptive Linear Neuron is a network having a single linear unit. It was developed by Widrow and Hoff in 1960. Some important points about Adaline are as follows

- It uses bipolar activation function.
- It uses delta rule for training to minimize the Mean-Squared Error (MSE) between the actual output and the desired/target output.
- The weights and the bias are adjustable.[7]

```
Algorithm 2: Adaline NN
Input: input matrix size ( \(4^{*} 3\) ), matrix wait ( \(\left.3^{*} 1\right)\) and desired matrix ( \(4^{*} 1\) )
Output: wait of size ( \(10 * 3\) )
Begin
Step 1 - Initialize the following to start the training
- Weights
- Bias
- Learning rate \(\alpha\)
For easy calculation and simplicity, weights and bias must be set to 0 and the learning rate must be set to 1 .
Step 2 - Continue step 3-8 when the stopping condition is not true.
Step 3 - Continue step 4-6 for every bipolar training pair s:t.
Step 4-Activate each input unit as follows -
\[
X_{i}=s_{i}(i=1 \text { to } n)
\]
```

Step 5 - Obtain the net input with the following relation -
$Y_{i n}=\mathrm{b}+\sum_{i}^{n} x_{i} * w_{i}$
Here ' $\mathbf{b}$ ' is bias and ' $\mathbf{n}$ ' is the total number of input neurons.
Step 6 - Apply the following activation function to obtain the final output -

$$
F\left(y_{i n}\right)= \begin{cases}1 & \text { if } \quad y_{i n} \geq 0 \\ -1 & \text { if } y_{\text {in }}<0\end{cases}
$$

Step 7 - Adjust the weight and bias as follows -
Case 1 - if $\mathbf{y} \neq \mathbf{t}$ then,
$w_{i}($ new $)=w_{i}($ old $)+\alpha\left(t-y_{i n}\right) x_{i}$
$b($ new $)=b($ old $)+\alpha\left(t-y_{i n}\right)$
Case 2 - if $\mathbf{y}=\mathbf{t}$ then,
$w_{i}($ new $)=w_{i}$ (old)
$b($ new $)=b($ old $)$
Here ' $\mathbf{y}$ ' is the actual output and ' t ' is the desired/target output.
( $t-y_{i n}$ ) is the computed error.
Step 8 - Test for the stopping condition, which will happen when there is no change in weight or the highest weight change occurred during training is smaller than the specified tolerance.
END
2.3 Random Number Generator Using Chaos Theory

Chaos theory is an area of mathematical science that studies the dynamic systems behavior; it is a good choice to use in designing of random number generator. In fact, chaotic system behaviors like; highly sensitive to initial states, mix up attribute, deterministic nature and also cannot predict the long term returns, so these properties are valuable in cryptography. Logistic equation of chaos which is explained below can be used to generate large pseudo random numbers, increase the randomness of produced keys and make the keys more secure [8].

### 2.3.1 Chaos Logistic Equation

It is a dynamic system second order deference equation; the standard form of the logistic equation is given by:
$\mathrm{F}(\mathrm{Yi})=\mathrm{R} \mathrm{Yi}(1-\mathrm{Yi})$
Where Yi is called the iteration of Y 0 (or population) and should by in subinterval $[0,1]$.
It all started to dawn on people when in 1960 a man named Edward Lorenz created a weather-model on his computer. Chaos is the study of complex systems with extreme sensitivity to slight changes in
conditions parameters. Chaos is the science of surprises, of the nonlinear and the unpredictable. It teaches us to expect the unexpected. While most traditional science deals with supposedly predictable phenomena like gravity, electricity, or chemical reactions, Chaos theory deals with nonlinear things that are effectively impossible to predict or control. [9].

```
Algorithm 3:1D logistic map chaos theory
Input:
- \(\mathrm{R}=3.6\)
- \(\quad \mathrm{y}=0.2\)
Output:
- binary string of length 5120 bits
Processing:
Begin
For \(\mathrm{i}=1\) to \(\mathbf{1 2 0}\)
Step1: Applying equation of 1d logistic map
Step2: Remove the negative signal
Step3: Cut 17 numbers after the comma
Step4: Save the numbers of each execution in buffer
Next for
Step5: Cuts final result to sub strings each of length 3, each sub string take mod 256; convert the result to
binary of 8 bits and aggregation the results in binary string
```

Step6: present final result of length 5120 bits
End

## 3. Proposed system design

Below the steps of the key generation method:
1: Input matrix size $(4 * 3)$, matrix wait ( $3 * 1$ ) and desired matrix $(4 * 1)$ to apply adaline neural network for 10 rounds only to get the wait of adaline.
2: Input matrix size $\left(3^{*} 3\right)$, applied Hopfield neural network on this matrix to present wait of Hopfield, The waiting from Hopfield can be broken or guessed so every number of waiting numbers is added with three serial numbers resulting from chaos theory (from the numbers collected from the chaos results saved in chaos buffer).
3: Create matrix $\mathrm{k}(10 * 3)$ by applying number of procedures will be carried out on the wait of adaline (step 1) .First, disposal from the negative signal. Second, cut two digits after comma. Third, apply the following equation to each number
$\mathrm{K}(\mathrm{i}, \mathrm{j})=(($ wait of adaline $(\mathrm{i}, \mathrm{j})+255) *(512 \wedge \mathrm{f}))$
Where $\mathrm{f}=3,4$, and 5
If number of column is 1 then $\mathrm{f}=3$
If number of column is 2 then $\mathrm{f}=4$
If number of column is 3 then $\mathrm{f}=5$
4: Create matrix $u(3 * 3)$ by applying number of procedures will be carried out on the wait of Hopfield (step 2).By apply the following equation to each number
$\mathrm{U}(\mathrm{i}, \mathrm{j})=\left(\right.$ wait of Hopfield $\left.(\mathrm{i}, \mathrm{j})+255^{\wedge} 6\right)$
5: Create an array $\mathbf{f}(10 * 3)$ by performing several operations between the last modification on the wait of Hopfield (matrix $\mathbf{u}(3,3)$ ) and the last modification on the wait of adaline (matrix $\mathbf{k}(10,3)$ )
First row of f equals from adding first row of Hopfield with first row of adaline.
Second row of f equals from adding second row of Hopfield with second row of adaline.
Third row of f equals from adding third row of Hopfield with third row of adaline.
Four row of $f$ equals from adding first column of Hopfield with fourth row of adaline.
Five row of $f$ equals from adding second column of Hopfield with five row of adaline.
Six row of $f$ is adding third column of Hopfield with six row of adaline.
Seven row of $f$ equals from adding Hopfield array $(1,0)$ with seven row of adaline.
Eight row of f equals from adding Hopfield array $(0,1)$ with eight row of adaline.
Nine row of f equals from adding Hopfield array $(1,2)$ with nine row of adaline.
Ten row of $f$ equals from adding Hopfield array $(2,1)$ with ten row of adaline.
6: Create an array $m\left(10^{*} 1\right)$ consist of
First row of $m$ equals to first row of $f$ array.
Second row of $m$ equals to second row of $f$ accumulated with the previous row.

Third row of $m$ equals to third row of $f$ accumulated with the previous row.
Four row of $m$ equals to four row of $f$ accumulated with the previous row.
Five row of $m$ array is adding five row of $f$ accumulated with the previous row.
Six row of $m$ equals to six row of $f$ accumulated with the previous row.
Seven row of $m$ equals to seven row of $f$ accumulated with the previous row.
Eight row of $m$ equals to eight row of $f$ accumulated with the previous row.
Nine row of $m$ equals to nine row of $f$ accumulated with the previous row.
Ten row of $m$ equals to ten row of $f$ accumulated with the previous row.
7: Convert each row of m to binary of length 64 bits so the final result is 640 bits
8: Making many operation:
8.1 Save the result of step 7 in buffer
8.2 Make DNA complement on the result of step 8.1 and add the result in buffer
8.3 Make initial permutation of DES algorithm on the result of step 8.2 and add it to the buffer
8.4 XOR operation between the results of steps 8.2 and 8.3 then add it to the buffer
8.5 Add with mode 256 between the results of steps 8.3 and 8.4 then add it to the buffer
8.6 Make initial permutation of DES algorithm on the result of step 8.5 then add it to the buffer
8.7 XOR operation between the results of steps 8.5 and 8.6 then add it to the buffer
8.8 Make DNA complement on the result of step 8.7 then add it to the buffer
-Now the length of binary string saved in buffer is 5120 bit
9: Apply chaos theory (1D logistic map) to present result of length 5120 bits
10: XOR operation between the result of step 8 and 9 to present the final key
11: Present final key of length 5120 bits
Example: - about (add with mode 256)
Suppose $\mathrm{a}=$ "11111111"
Suppose b="00000111"

## Solve:

Convert a to integer=255 and Convert b to integer=7
$=(255+7) \bmod 256=262 \bmod 256=6$
Convert the result to binary="00000110"

-Now the length of binary string saved in buffer is 5120 bit
Step 9: Apply chaos theory (1D logistic map) to give result of length 5120 bits
Step 10: XOR operation between the result of step 8 and 9 to present the final key
Step 11: Present the final key of 5120 bits
End


## Example

Adaline is given the four different input and output combinations $\quad W 0=-0.12, \quad w 1=0.4, \quad w 2=0.65$

| $X_{0}$ | $X_{1}$ | $X_{2}$ | $Y_{\text {desired }}$ |
| :--- | :--- | :--- | :--- |
| 1 | -1 | -1 | -1 |
| 1 | -1 | 1 | -1 |
| 1 | 1 | -1 | 1 |
| 1 | 1 | 1 | -1 |

1. First the input pattern: $+1 \quad-1 \quad-1$

Weights: -0.12 $0.4 \quad 0.65$

$$
\begin{aligned}
& \text { net }=\sum_{i=1}^{n} \mathrm{X}_{\mathrm{i}} \cdot \mathrm{~W}_{\mathrm{i}} \\
&=(+1 *-0.12)+(-1 * 0.4)+(-1 * 0.65)=-1.17 \\
&\mathrm{~d}=\text { desired output }=-1 \text { (for first pattern })
\end{aligned}
$$

$\therefore \delta=\mathrm{d}-$ net
$=-1-(-1.17)=0.17$
$\eta=0.1$
Also we must compute:-
$\Delta \mathrm{w}_{\mathrm{ij}}=\eta \cdot \delta \cdot \mathrm{x}_{\mathrm{i}}$
For convenience, these figures have been rounded to two places after the decimal point, so become:- $\eta \cdot \delta \cdot x_{0}=(0.1 * 0.17 *+1)=0.017 \approx=0.02$
$\eta \square \mathrm{x}_{1}=0.1^{*} 0.17 *(-1)=-0.017 \approx-0.02$
$\eta \square \mathrm{x}_{2}=0.1 * 0.17 *(-1)=-0.017 \approx-0.02$
Note: - We continue to work with the rest of the input and thus we get the following results

| $\Delta w_{0} \Delta w_{1} \Delta w_{2}$ |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $X_{0}$ | $X_{1}$ | $X_{2}$ | $W_{0}$ | $W_{1}$ | $W_{2}$ | Net | d | $\eta \delta \mathrm{x}_{0}$ | $\eta \delta \mathrm{x}_{1}$ | $\eta \delta \mathrm{x}_{2}$ |
| +1 | -1 | -1 | -0.12 | 0.40 | 0.65 | -1.17 | -1 | 0.02 | -0.02 | -0.02 |
| +1 | -1 | +1 | -0.12 | 0.40 | 0.65 | 0.13 | -1 | -0.11 | 0.11 | -0.11 |
| +1 | +1 | -1 | -0.12 | 0.40 | 0.65 | -0.37 | +1 | 0.14 | 0.14 | -0.14 |
| +1 | +1 | +1 | -0.12 | 0.40 | 0.65 | 0.93 | -1 | -0.19 | -0.19 | -0.19 |

meanj $=\operatorname{total}(\Delta w i j) / p \quad \mathrm{p}=4$
Mean $_{0}=-0.14 / 4=-0.035=-0.04$
Mean $_{1}=-0.04 / 4=-0.01$
Mean $_{2}=-0.46 / 4=-0.115=-0.12$
$\therefore \mathrm{W}_{\mathrm{ij}}{ }^{\text {new }}=\mathrm{W}_{\mathrm{ij}}^{\text {old }}+$ meanj
$W_{0}^{\text {new }}=-0.12+(-0.04)=-0.16$
$\mathrm{W}_{1}^{\text {new }}=0.40+(0.01)=-0.41$
$\mathrm{W}_{2}{ }^{\text {new }}=0.65+(-0.12)=0.53$
Continue until 10 round

| No. of round | W0 | W1 | W2 |
| :---: | :---: | :---: | :---: |
| Round1 | -0.14 | 0.04 | -0.46 |
| Round2 | -0.14 | 0.06 | -0.46 |
| Round3 | -0.13 | 0.09 | -0.49 |
| Round4 | -0.13 | 0.11 | -0.49 |
| Round5 | -0.11 | 0.13 | -0.51 |
| Round6 | -0.11 | 0.17 | -0.53 |
| Round7 | -0.1 | 0.18 | -0.54 |
| Round8 | -0.09 | 0.21 | -0.57 |
| Round9 | -0.09 | 0.25 | -0.59 |
| Round10 | -0.07 | 0.27 | -0.61 |

## 2. Consider the following samples are inputs for Hopfield:



1-use Hebb rule to find the weights matrix

$$
\begin{array}{lll}
\mathrm{w}_{00} & \mathrm{w}_{01} & \mathrm{w}_{02} \\
\mathrm{w}_{10} & \mathrm{w}_{11} & \mathrm{w}_{12} \\
\mathrm{w}_{20} & \mathrm{w}_{21} & \mathrm{w}_{22} \\
\\
\mathrm{w}_{00}=\mathrm{w}_{11}=\mathrm{w}_{22}=0 \text { because } \mathrm{i}=\mathrm{j} \\
\mathrm{w}_{01}=(1 * 1)+(1 *-1)=(-1 * 1)=-1=\mathrm{w}_{10} \\
\mathrm{w}_{02}=(1 * 1)+(1 * 1)=(-1 * 1)=1=\mathrm{w}_{20} \\
\mathrm{w}_{12}=(1 * 1)+(-1 * 1)=(1 * 1)=1=\mathrm{w}_{21}
\end{array}
$$

$\mathrm{w}=\left[\begin{array}{ccc}0 & -1 & 1 \\ -1 & 0 & 1 \\ 1 & 1 & 0\end{array}\right]$
The waiting from Hopfield can be broken or guessed so every number of waiting numbers is add with three serial numbers resulting from chaos theory (from the numbers collected from the chaos results after the elimination of the negative signal and cut 17 rank after the comma)
The number produced from chaos is: $576879206438232902231654485 \ldots \ldots$.
The new wait is:
$W_{11}=0+576=576$
$W_{12}=-1+879=878$ and so on
New $w=\left[\begin{array}{ccc}576 & 878 & 207 \\ 437 & 232 & 903 \\ 232 & 655 & 485\end{array}\right]$

## 3. Update the wait of adaline

$\mathrm{W}_{00}=(14+255) *\left(512^{\wedge} 3\right)=269 * 134217728=36104568832$
$\mathrm{W}_{10}=(14+255) *\left(512^{\wedge} 3\right)=269 * 134217728=36104568832$
$\mathrm{W}_{20}=(13+255) *\left(512^{\wedge} 3\right)=268 * 134217728=35970351104$
$\mathrm{W}_{30}=(13+255) *(512 \wedge 3)=268 * 134217728=35970351104$
$\mathrm{W}_{40}=(11+255) *\left(512^{\wedge} 3\right)=266^{*} 134217728=35701915648$
$\mathrm{W}_{50}=(11+255) *\left(512^{\wedge} 3\right)=266^{*} 134217728=35701915648$
$\mathrm{W}_{60}=(10+255) *\left(512^{\wedge} 3\right)=256 * 134217728=35567697920$

```
W
W
W}\mp@subsup{}{90}{}=(7+255)*(512^3) =262*134217728= 3516504473
W}\mp@subsup{0}{01}{}=(4+255)*(512^4) =259*68719476736= 17798344474624
W}\mp@subsup{W}{11}{}=(6+255)*(512^4)=261*68719476736= 17935783428096
W}\mp@subsup{2}{21}{}=(9+255)*(512^4)=264*68719476736= 18141941858304
W}\mp@subsup{W}{31}{}=(11+255)*(512^4)=266*68719476736= 18279380811776
W}\mp@subsup{W}{41}{}=(13+255)*(512^4)=268*68719476736=1841681976524
W}\mp@subsup{\mp@code{51}}{}{\prime}=(17+255)*(512^4)=272*68719476736= 18691697672192
W
W
W
W}\mp@subsup{\textrm{O}}{91}{}=(27+255)*(51\mp@subsup{2}{}{\wedge}4)=282*68719476736= 1937889243955
W}\mp@subsup{\textrm{O}}{02}{}=(46+255)*(512^5) =301*35184372088832= 10590495998738432
W}\mp@subsup{W}{12}{}=(46+255)*(512^5) =301*35184372088832= 10590495998738432
W}\mp@subsup{W}{22}{}=(49+255)*(512^5) =304*35184372088832= 10696049115004928
W W32 = (49+255)*(512^5) =304*35184372088832= 10696049115004928
W}\mp@subsup{W}{42}{}=(51+255)*(512^5)=306*35184372088832=10766417859182592
W
W}\mp@subsup{W}{62}{}=(54+255)*(512^5)=309*35184372088832=10871970975449088
W}\mp@subsup{\textrm{T}}{72}{}=(57+255)*(512^5)=312*35184372088832=1097752409171558
W
W}\mp@subsup{\mp@code{92}}{}{=}(61+255)*(512^5)=316*35184372088832= 11118261580070912
```


## 4. Update Hopfield waits

$\mathrm{W}_{00}=\mathrm{W}_{00} *\left(255^{\wedge} 6\right)=576 *(274941996890625)=158366590209000000$
$\left.\mathrm{W}_{01}=\mathrm{W}_{01} *(255 \wedge 6)=878 *(274941996890625)\right)=241399073269968750$
$\left.\mathrm{W}_{02}=\mathrm{W}_{02} *\left(255^{\wedge} 6\right)=207 *(274941996890625)\right)=56912993356359375$
$\left.\mathrm{W}_{10}=\mathrm{W}_{10} *\left(255^{\wedge} 6\right)=437 *(274941996890625)\right)=120149652641203125$
$\mathrm{W}_{11}=\mathrm{W}_{11} *(255 \wedge 6)=232 *(274941996890625)=63786543278625000$
$\mathrm{W}_{12}=\mathrm{W}_{12} *(255 \wedge 6)=903 *(274941996890625)=248272623192234375$
$\mathrm{W}_{30}=\mathrm{W}_{30} *\left(255^{\wedge} 6\right)=232 *(274941996890625)=63786543278625000$
$\mathrm{W}_{31}=\mathrm{W}_{31} *\left(255{ }^{\wedge} 6\right)=655 *(274941996890625)=180087007963359375$
$\mathrm{W}_{32}=\mathrm{W}_{32} *(255 \wedge 6)=485 *(274941996890625)=133346868491953125$

## 5. Create matrix $(10,3)$ include many ope ration between modified wait of adaline and modified wait of Hopfield

$\mathrm{W}_{00}=36104568832+158366590209000000=158366626313568832$
$\mathrm{W}_{01}=17798344474624+241399073269968750=241416871614443374$
$\mathrm{W}_{02}=10590495998738432+56912993356359375=67503489355097807$
$\mathrm{W}_{10}=36104568832+120149652641203125=120149688745771957$
$\mathrm{W}_{11}=17935783428096+63786543278625000=63804479062053096$
$\mathrm{W}_{12}=10590495998738432+248272623192234375=258863119190972807$
$\mathrm{W}_{21}=35970351104+63786543278625000=63786579248976104$
$W_{22}=18141941858304+180087007963359375=180105149905217679$
$W_{23}=10696049115004928+133346868491953125=144042917606958053$
$\mathrm{W}_{30}=35970351104+158366590209000000=158366626179351104$
$W_{31}=18279380811776+120149652641203125=120167932022014901$
$\mathrm{W}_{32}=10696049115004928+63786543278625000=74482592393629928$
$W_{40}=35701915648+241399073269968750=241399108971884398$
$W_{41}=18416819765248+63786543278625000=63804960098390248$
$W_{42}=10766417859182592+180087007963359375=190853425822541967$
$\mathrm{W}_{50}=35701915648+56912993356359375=56913029058275023$
$\mathrm{W}_{51}=18691697672192+248272623192234375=248291314889906567$
$W_{52}=10836786603360256+133346868491953125=144183655095313381$
$\mathrm{W}_{60}=35567697920+120149652641203125=120149688208901045$
$\mathrm{W}_{61}=18760417148928+120149652641203125=120168413058352053$
$\mathrm{W}_{62}=10871970975449088+120149652641203125=131021623616652213$
$\mathrm{W}_{70}=35433480192+241399073269968750=241399108703448942$
$\mathrm{W}_{71}=18966575579136+241399073269968750=241418039845547886$
$\mathrm{W}_{72}=10977524091715584+241399073269968750=252376597361684334$
$\mathrm{W}_{80}=35433480192+248272623192234375=248272658625714567$
$\mathrm{W}_{81}=19241453486080+248272623192234375=248291864645720455$
$\mathrm{W}_{82}=11047892835893248+248272623192234375=259320516028127623$
$\mathrm{W}_{90}=35165044736+180087007963359375=180087043128404111$
$\mathrm{W}_{91}=19378892439552+180087007963359375=180106386855798927$
$\mathrm{W}_{92}=11118261580070912+180087007963359375=191205269543430287$
6. Create matrix $\mathbf{m}(10,1)$ include many ope ration based on previous matrix
$\mathrm{m}_{00}=\mathrm{w}_{00}+\mathrm{w}_{01}+\mathrm{w}_{02}=158366626313568832+241416871614443374$
$+67503489355097807=467286987283110013$
$\mathrm{m} 10=\mathrm{w}_{10}+\mathrm{w}_{11}+\mathrm{w}_{12}+\mathrm{m}_{00}=910104274281907873$
$\mathrm{m}_{20}=\mathrm{w}_{20}+\mathrm{w}_{21}+\mathrm{w}_{22}+\mathrm{m}_{10}=1298038921043059709$
$\mathrm{m}_{30}=\mathrm{w}_{30}+\mathrm{w}_{31}+\mathrm{w}_{32}+\mathrm{m}_{20}=1651056071638055642$
$\mathrm{m}_{40}=\mathrm{w}_{40}+\mathrm{w}_{41}+\mathrm{w}_{42}+\mathrm{m}_{30}=2147113566530872255$
$\mathrm{m}_{50}=\mathrm{w}_{50}+\mathrm{w}_{51}+\mathrm{w}_{52}+\mathrm{m}_{40}=2596501565574367226$
$\mathrm{m}_{60}=\mathrm{w}_{60}+\mathrm{w}_{61}+\mathrm{w}_{62}+\mathrm{m}_{50}=2967841290458272537$
$\mathrm{m}_{70}=\mathrm{w}_{70}+\mathrm{w}_{71}+\mathrm{w}_{72}+\mathrm{m}_{60}=3703035036368953699$
$\mathrm{m}_{80}=\mathrm{w}_{80}+\mathrm{w}_{81}+\mathrm{w}_{82}+\mathrm{m}_{70}=4458920075668516344$
$\mathrm{m}_{90}=\mathrm{w}_{90}+\mathrm{w}_{91}+\mathrm{w}_{92}+\mathrm{m}_{80}=5010318775196149669$

## 7. Applying many operations such that:

7.1-Convert each of $\left(m_{00}, m_{10}, m_{20}, m_{30}, m_{40}, m_{50}, m_{60}, m_{70}, m_{80}, m_{90}\right)$ to binary of 64 bits, the result is: 000001100111110000100011000010100000010100.
7.2- Applying DNA complement on the result of step 7.1, the result is:

111110011000001111011100111101011111101011
7.3-Applying initial permutation of 64 bits on the result of step 7.2 , the result is:-
0111110100111101010011000110101111111111010.
7.4-Applying XOR operation between step 7.2 and step 7.3, the result is:
$1000010010111110100100001001111000000101100 \ldots . .$.
7.5-Applying add with mod 256 between result of step 7.3 and 7.4 , the result is:-

00000001111110111101110000001001000001001110
7.6-applying initial permutation of 64 bits on the result of step 7.5 , the result is:-
$10100110000001100101010001101011001001101010001 \ldots$.
7.7-applying XOR operation between step 7.5 and 7.6, the result is:-
1010011111111101100010000110001000100010010000......
7.8-applying DNA complement on the result of step 7.7, the result is: 01011000000000100111011110011101110111011011111.
8. Aggregation the results of steps (7.1, 7.2,... 7.8) to present the binary string of 5120 bit
9. Applying chaos method (1D Logistic map), the result is:

0100000000000000101011110110111100011000100110001100111
10. Applying XOR operation between step 8 and 9, the result is:-
010001100111110010001100011001010001110110110100110. $\qquad$
11. Present key of length 5120 bits
12. Implementation and evaluation result


Figure 3-Result of Hopfield and adaline
This figure appears wait of adaline, wait of Hopfield, modified adaline wait, modified Hopfield wait and combine between two waits to generate matrix $\mathrm{m}\left(10^{*} 1\right)$


Figure 4- final key
This figure appears the final result of key generation in binary of length 5120 bits, after applied many operation in order to access the highest random key

Based on the results of the main 5 statistic tests which have been performed on (Towards generation robust key based on neural networks and chaos theory), as shown in table (1). The suggested modified method showed better results, particularly on the test of frequency, considered as the most important one of the five tests.

Table 1-Results of the 5 Tests that have been applied to the proposed key ( 5120 bits)

| Statistical Test |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Tests |  | Freedom Degree | First sample | Second sample | Third Sample |
| Frequency Test |  | Must be $<=3.84$ | Pass $=0.013$ | Pass $=1.188$ | Pass $=2.813$ |
| Run <br> Test | T0 | $\begin{aligned} & \text { Must be } \\ & <=19.391 \end{aligned}$ | Pass $=3.856$ | Pass $=11.731$ | Pass $=11.450$ |
|  | T1 |  | Pass $=5.622$ | Pass $=11.686$ | Pass $=9.264$ |
| Poker Test |  | $\begin{gathered} \hline \text { Must be } \\ <=11.1 \end{gathered}$ | Pass $=2.572$ | Pass $=2.659$ | Pass $=5.975$ |
| Serial Test |  | $\begin{gathered} \hline \text { Must be } \\ <=7.81 \end{gathered}$ | Pass $=0.219$ | Pass $=5.034$ | Pass $=3.269$ |
| Auto Correlation Test | $\begin{aligned} & \hline \text { Shift } \\ & \text { No. } 1 \end{aligned}$ | Must be$<=3.84$ | Pass $=0.086$ | Pass $=0.877$ | Pass $=0.267$ |
|  | Shift <br> No. 2 |  | Pass $=0.028$ | Pass $=0.113$ | Pass $=1.070$ |
|  | Shift <br> No. 3 |  | Pass $=2.408$ | Pass $=0.016$ | Pass $=1.618$ |
|  | Shift <br> No. 4 |  | Pass $=1.955$ | Pass $=3.831$ | Pass $=0.704$ |
|  | Shift <br> No. 5 |  | Pass $=1.159$ | Pass $=1.220$ | Pass $=0.057$ |
|  | Shift <br> No. 6 |  | Pass $=1.190$ | Pass $=1.251$ | Pass $=0.050$ |
|  | Shift <br> No. 7 |  | Pass $=0.728$ | Pass $=2.863$ | Pass $=0.071$ |
|  | Shift <br> No. 8 |  | Pass $=1.014$ | Pass $=0.153$ | Pass $=1.515$ |
|  | $\begin{gathered} \hline \text { Shift } \\ \text { No. } 9 \end{gathered}$ |  | Pass $=0.240$ | Pass $=1.284$ | Pass $=0.002$ |
|  | $\begin{gathered} \hline \text { Shift } \\ \text { No. } 10 \end{gathered}$ |  | Pass $=0.132$ | Pass $=0.003$ | Pass $=0.013$ |

## 13. Conclusions

1. When the initial values of the chaos theory change, all the results change according to the initial values
2. It is impossible to think of the initial values of the chaos theory because its number of probability is too large
3. Each step of the key generation depends on the results of the previous step and this has a big role in the difficulty of breaking the key
4. It was concluded that neural networks did not give random results because they are trained so a chaos equation was used to generate clutter numbers to increase clutter in the key and to form a random and more secure key.
5. Artificial intelligence algorithms have been used with chaos in the producing of keys and are one of the new ways that not mentioned before.

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[^0]:    * Email: 110032@uotechnology.edu.iq

